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Technical Reference Derivation of the Shannon Spectral Efficiency Limit 1 (4) Time Rev No. Reference 10-Apr-2014 PA2 n/a 22:22 shannonlimit.tex

#### Derivation of the Shannon Spectral Efficiency Limit

Date

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### 1 Introduction

In digital communication systems, spectral efficiency,  $\eta$ , is defined as

$$\eta = \frac{C}{W},\tag{1}$$

where *C* is the channel capacity (in [bits/second]) and *W* is the channel bandwidth (in [Hz]).

Spectral efficiency provides a measure of how efficiently a communication scheme utilizes bandwidth. For any fixed capacity *C*, spectral efficiency approaches zero as the bandwidth approaches infinity. This is the absolute worst efficiency possible (infinite bandwidth!).

Shannon's famous capacity theorem can be rearranged to provide a relationship between spectral efficiency and SNR, or  $E_b/N_0^{-1}$  (see Figure 1). It appears from the plot that the relationship approaches an asymptotic limit as  $\eta \rightarrow 0$ . In this paper we show this limit is ln 2, or approximately -1.59 dB.

### 2 Derivation

Shannon's capacity theorem is stated in [2] as follows:

**Shannon's Capacity Theorem.** Let P be the average transmitter power, and suppose the noise is white thermal noise of power N in the band W. By sufficiently complicated encoding systems it is possible to transmit binary digits at a rate

$$C = W \log_2 \left( 1 + \frac{P}{N} \right), \tag{2}$$

with as small a frequency of errors as desired. It is not possible by any encoding method to send at a higher rate and have an arbitrarily low frequency of errors.

If the energy per bit transmitted is denoted  $E_b$  and we are transmitting *C* bits per second, then the signal power  $P = CE_b$ . Since the noise is white, the noise power in a bandwidth *W* is  $WN_0$  where  $N_0$  is the noise spectral density, and we may rewrite the capacity equation as

$$\frac{C}{W} = \log_2\left(1 + \frac{E_b}{N_0}\frac{C}{W}\right).$$
(3)

<sup>&</sup>lt;sup>1</sup>For a good discussion on SNR and  $E_b/N_0$  see section 5.2.2 in [1].

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If we let  $\eta = C/W$  (the spectral efficiency), then we can reexpress in terms of  $E_b/N_0$  as

$$\frac{E_b}{N_0} = \frac{2^{\eta} - 1}{\eta},$$
(4)

that is, we can express  $E_b/N_0$  as a function of  $\eta$ , which is plotted in Figure 1.

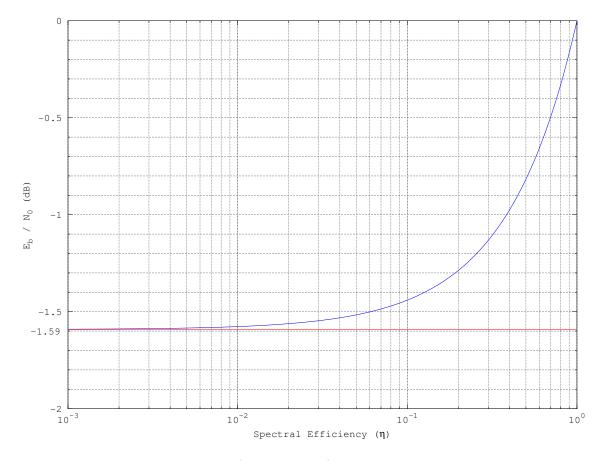


Figure 1: Capacity Curve

The red line is the asymptote of the curve as  $\eta$  approaches zero and is approximately -1.59 dB (ln 2). We derive this result as follows.

The asymptotic value we seek, *a*, is the limit of  $E_b/N_o$  as  $\eta$  approachs zero.

$$a = \lim_{\eta \to 0} \frac{2^{\eta} - 1}{\eta}.$$
(5)

Since the argument of the limit is a rational function  $f(\eta)/g(\eta)$  and f(0) = g(0) = 0, we may apply L'Hopital's rule, which states [3] that, under these conditions,

$$\lim_{\eta \to 0} \frac{f(\eta)}{g(\eta)} = \lim_{\eta \to 0} \frac{f'(\eta)}{g'(\eta)}.$$
(6)

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In order to take the derivative of the numerator, express the value 2 in terms of the natural logarithm:

$$2 = e^{\ln 2}.\tag{7}$$

Then

$$2^{\eta} = \left(e^{\ln 2}\right)^{\eta} \tag{8}$$

$$=e^{\eta\ln 2} \tag{9}$$

and

$$\frac{\mathrm{d}2^{\eta}}{\mathrm{d}\eta} = \frac{\mathrm{d}e^{u}}{\mathrm{d}u}\frac{\mathrm{d}u}{\mathrm{d}\eta},\tag{10}$$

where  $u = \eta \ln 2$ . Therefore

$$a = \lim_{\eta \to 0} \frac{\mathrm{d}2^{\eta}}{\mathrm{d}\eta} \tag{11}$$

$$=\lim_{\eta\to 0} e^u \ln 2 \tag{12}$$

$$= \lim_{\eta \to 0} \left( e^{\ln 2} \right)^{\eta} \ln 2$$
 (13)

$$= \ln 2. \blacksquare \tag{14}$$

#### 3 Conclusion

In a digital communication system, we can trade off bandwidth (spectral efficiency) for  $E_b/N_0$ , as seen from Figure 1. For example, a spectral efficiency of one requires an  $E_b/N_0$  of approximately 0 dB to reach capacity. But if we allow our spectral efficiency to be reduced to 0.1, we can reach the Shannon capacity limit with 1.45 dB less power.

However, this tradeoff has a limit in that, no matter how much bandwidth you allow a signal to have, you can never use less than -1.59 dB  $E_b/N_0$  and still achieve capacity.

## References

- [1] Bernard Sklar, Digital Communications, 2nd ed. Prentice Hall P T R, 2001.
- [2] C. E. Shannon, "Communication in the presence of noise," *Proceedings of the Institute of Radio Engineers*, vol. 37, pp. 10–21, 1949.
- [3] S. K. Stein, Calculus and Analytic Geometry, 2nd ed. McGraw-Hill, 1973.